## Solution to Matching Problems

Exercise 1 Construct an example in which there is more than one stable matching. (You only need two boys and two girls to do this.)

Solution 1 Suppose the preferences are:
Alan: Megan $\succ$ Melissa.
Ron: Melissa $\succ$ Megan.
Melissa: Alan $\succ$ Ron.
Megan: Ron $\succ$ Alan.
So the match Alan-Megan, Ron-Melissa is stable. (This is the favorite of the boys.) The match Alan-Melissa, Ron-Megan is also stable. (This is the favorite of the girls.)

Exercise 2 Suppose that the boys all have different favorite girls. How many steps does it take for the algorithm to converge?

Solution 2 One. They all propose to a different girl.
Exercise 3 Suppose that the boys have identical preferences. How many steps does it take for the algorithm to converge?

Solution 3 At the first round, all boys propose to the same girl. She "dates" her favorite. At the second round, all but that favorite one boy proposes to the second-best girl. She "dates" her favorite. The process continues for $N$ rounds until one guy is left to propose to the least desirable girl.

Exercise 4 Suppose preferences are given by the following tables:

| BOY | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Adam | Beth | Amy | Diane | Ellen | Cara |
| Bill | Diane | Beth | Amy | Cara | Ellen |
| Carl | Beth | Ellen | Cara | Diane | Amy |
| Dan | Amy | Diane | Cara | Beth | Ellen |
| Eric | Beth | Diane | Amy | Ellen | Cara |

Boys' Preferences

| GIRL | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Amy | Eric | Adam | Bill | Dan | Carl |
| Beth | Carl | Bill | Dan | Adam | Eric |
| Cara | Bill | Carl | Dan | Eric | Adam |
| Diane | Adam | Eric | Dan | Carl | Bill |
| Ellen | Dan | Bill | Eric | Carl | Adam |

Girls' Preferences
Find a stable matching using the Gale-Shapley algorithm with boys making proposals. Find a stable matching using the Gale-Shapley algorithm with girls making proposals.

Solution 4 If the boys make proposals, then you get:
Round 1:

| GIRL | Date | Rejected this round |
| :---: | :---: | :---: |
| Amy | Dan |  |
| Beth | Carl | Adam, Eric |
| Cara |  |  |
| Diane | Bill |  |
| Ellen |  |  |

Round 2:

| GIRL | Date | Rejected this round |
| :---: | :---: | :---: |
| Amy | Adam | Dan |
| Beth | Carl |  |
| Cara |  |  |
| Diane | Eric | Bill |
| Ellen |  |  |

Round 3:

| GIRL | Date | Rejected this round |
| :---: | :---: | :---: |
| Amy | Adam |  |
| Beth | Carl | Bill |
| Cara |  |  |
| Diane | Eric | Dan |
| Ellen |  |  |

Round 4:

| GIRL | Date | Rejected this round |
| :---: | :---: | :---: |
| Amy | Adam |  |
| Beth | Carl |  |
| Cara | Dan | Bill |
| Diane | Eric |  |
| Ellen |  |  |

Round 5:

| GIRL | Date | Rejected this round |
| :---: | :---: | :---: |
| Amy | Adam | Bill |
| Beth | Carl |  |
| Cara | Dan |  |
| Diane | Eric |  |
| Ellen |  |  |

Round 6:

| GIRL | Date | Rejected this round |
| :---: | :---: | :---: |
| Amy | Adam |  |
| Beth | Carl |  |
| Cara | Bill | Dan |
| Diane | Eric |  |
| Ellen |  |  |

Round 7:

| GIRL | Date | Rejected this round |
| :---: | :---: | :---: |
| Amy | Adam |  |
| Beth | Carl | Dan |
| Cara | Bill |  |
| Diane | Eric |  |
| Ellen |  |  |

Round 8:

| GIRL | Date | Rejected this round |
| :---: | :---: | :---: |
| Amy | Adam |  |
| Beth | Carl |  |
| Cara | Bill |  |
| Diane | Eric |  |
| Ellen | Dan |  |

This is the boy-optimal stable matching. If girls make proposals, then we have:

Round 1:

| Boy | Date | Rejected this round |
| :---: | :---: | :---: |
| Adam | Diane |  |
| Bill | Cara |  |
| Carl | Beth |  |
| Dan | Ellen |  |
| Eric | Amy |  |

This is the girl-optimal stable matching.
Exercise 5 This exercise shows that stable matchings need not exist if there are not "two sides." Consider the following "roommate" problem. There are four people, Pat, Chris, Dana, and Leslie. They must pair off (each pair will share a two-bed suite). Each has preferences over which of the others they would prefer to have as a roommate. The preferences are:

Leslie: Pat $\succ$ Chris $\succ$ Dana
Chris: Leslie $\succ$ Pat $\succ$ Dana
Pat: Chris $\succ$ Leslie $\succ$ Dana
Dana: Chris $\succ$ Leslie $\succ$ Pat
Show that no stable matching exists. (That is, no matter who you put together, they will always be two potential roommates who are not matched, but prefer each other to their current roommate.)

Solution 5 Someone must be matched with Dana. Whoever that is would prefer either of the other two to Dana. One the other hand, one of the others will think that Dana's partner is the best roommate. (Concretely, if the pairings are Dana-Chris and Leslie-Pat are paired, then Pat and Chris would prefer to be with each other than their mate. The same kind of argument would work whoever paired with Dana.)

